

# Algebraic model structures

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18 September, 2011

Homotopy Theory and Its Applications

AWM Anniversary Conference

ICERM

- 1 Flavors of the theory of algebraic model structures
- 2 The monoidal algebraic model structure on **sSet**

# Weak factorization systems

Any model structure has two **weak factorization systems** (**wfs**):

- (cofibrations, trivial fibrations)
- (trivial cofibrations, fibrations)

The left and right classes satisfy **factorization** and **lifting** axioms.

## Examples

- (retracts of rel. cell complexes, trivial Serre fibrations) on **Top**
- (anodyne extensions, Kan fibrations) on **sSet**
- (injections with projective cokernel, surjections) on  **$A$ -mod**
- (monomorphisms, epimorphisms) on **Set**

# “Algebraic” perspective

## Thinking “algebraically”

to characterize maps or objects satisfying a certain *property*, assign to each one a particular *structure* that demonstrates the property.

## Examples

- a surjective map in **Set** or  **$A$ -mod** admits a (set-based) **section**
- a relative cell complex admits a **cellular decomposition**
- for any Kan complex, can choose **fillers for all horns**

## (Co)algebras for (co)monads

For all of these examples there is a **monad** or a **comonad** whose **algebras** or **coalgebras** have exactly this form.

# Can homotopy theory be made algebraic?

Answer: yes!

Cofibrantly generated model categories admit **algebraic model structures**:

- a fibrant-replacement monad and a cofibrant-replacement comonad
- fibrations and trivial fibrations are algebras for a pair of monads
- **cellular** cofibrations and trivial cofibrations are coalgebras
- the (co)monads define the functorial factorizations
- the (co)algebra structures give explicit solutions to lifting problems

## Algebraic weak factorization systems

In place of wfs, we use **algebraic weak factorization systems (awfs)**:

- factorization defines a monad and a comonad on the arrow category
- maps in the right class admit pre-**algebra** structures
- maps in the left class admit pre-**coalgebra** structures

# A classical example: Hurewicz fibrations

Consider the wfs (cofibrations and htpy equiv, fibrations) on **Top**

- A map  $f: E \rightarrow B$  can be factored through the space of **Moore paths**

$$E \xrightarrow{\iota f} \Gamma_B E \xrightarrow{\pi f} B$$

- $f \mapsto \pi f$  is a monad on **Top**/ $B$ , or better: on **Top**<sup>2</sup>
- pre-algebras are maps with path lifting functions; ie Hurewicz fibrations
- $f \mapsto \iota f$  is a comonad on **Top**<sup>2</sup>; coalgebras are maps in the left class
- algebra and coalgebra structures can be used to construct lifts

$$\begin{array}{ccc} A & \longrightarrow & E \\ \downarrow \sim & \nearrow & \downarrow \\ X & \longrightarrow & B \end{array}$$

# Cellularity: definition

## Cofibrantly generated awfs

An awfs is **cofibrantly generated** if there exists a set  $\mathcal{I}$  of arrows such that the right class equals those maps that lift against  $\mathcal{I}$ .

## Baby example

In **Set**,  $\mathcal{I} = \{\emptyset \rightarrow *\}$  generates (monomorphism, epimorphism).

## Lemma (R.)

In a cofibrantly generated awfs, all right maps admit algebra structures.

## Cellular maps

A map (in the left class) is **cellular** if it admits a coalgebra structure.

# Cellularity: examples

## Examples

category	generators	cofibrations	cellular cofibrations
<b><math>A</math>-mod</b>	$\{0 \rightarrow A\}$	monos w/ projective cokernel	monos w/ free cokernel
<b>Top</b>	$\{S^{n-1} \rightarrow D^n\}$	retracts of relative cell cxes	relative cell cxes
<b>sSet</b>	$\{\partial\Delta^n \rightarrow \Delta^n\}$	monomorphisms	monomorphisms
<b>sSet</b>	$\{\Lambda_k^n \rightarrow \Delta^n\}$	anodyne extensions	"anodyne cell cxes"?



# Algebraic Quillen adjunctions by example

## Sample Theorem (R.)

$| - | : \mathbf{sSet} \rightleftarrows \mathbf{Top} : S$  is an **algebraic Quillen adjunction**.

- all cofibrations in **sSet** are **cellular**, filtered by attaching stages
- images under  $| - |$  not just cofibrations but **cellular**—here, relative cell complexes—with a specified **algebraic structure**—here, a cellular decomposition
- **algebraic Serre fibrations** are equipped with chosen lifted homotopies;

$$\begin{array}{ccc} |\Lambda_k^{n+1}| \cong D^n & \longrightarrow & X \\ \downarrow & \searrow & \downarrow f \\ |\Delta^{n+1}| \cong D^n \times I & \longrightarrow & Y \end{array} \quad \mapsto \quad \begin{array}{ccc} \Lambda_k^{n+1} & \longrightarrow & SX \\ \downarrow & \searrow & \downarrow Sf \\ \Delta^{n+1} & \longrightarrow & SY \end{array}$$

- images under  $S$  are **algebraic Kan fibrations** with chosen horn fillers

# Existence of algebraic Quillen adjunctions

## In an algebraic Quillen adjunction

the left adjoint lifts to commuting functors of coalgebras and the right adjoint lifts to commuting functors of algebras.

Modulo the usual acyclicity condition and a compatibility condition which is not the main point:

## Cellularity Theorem (R.) Cellularity & Uniqueness Theorem (R.)

Suppose  $\mathcal{M}$  has an algebraic model structure generated by  $\mathcal{J}$  and  $\mathcal{I}$ ,  $\mathcal{K}$  has an algebraic model structure, and  $F: \mathcal{M} \rightleftarrows \mathcal{K}: U$ . Then  $F \dashv U$  is an **algebraic Quillen adjunction** iff  $F\mathcal{J}$  and  $F\mathcal{I}$  are **cellular**. Furthermore, the coalgebra structures assigned to  $F\mathcal{I}$  and  $F\mathcal{J}$  determine everything.

## Corollary (R.)

Whenever an algebraic model structure is lifted along an adjunction, the adjunction is canonically an algebraic Quillen adjunction.

# The monoidal model structure on $\mathbf{sSet}$

The combinatorics necessary to prove theorems such as

## Theorem (Quillen?)

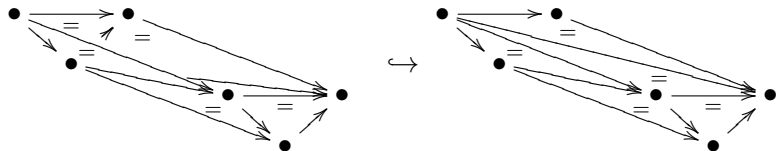
If  $X$  is a Kan complex and  $A$  is a simplicial set then  $X^A$  is a Kan complex.

are encoded in the fact that  $\mathbf{sSet}$  is a **monoidal model category**. Precisely:

## Equivalent Theorem (pushout-product axiom)

The pushout-product of an anodyne extension with a monomorphism is an anodyne extension.

eg,  $(\Lambda_1^2 \rightarrow \Delta^2) \hat{\times} (\partial\Delta^1 \rightarrow \Delta^1)$  is



# Two-variable adjunctions

The equivalence is because both theorems describe the interaction between wfs and a **two-variable adjunction**.

## Definition

A **two-variable adjunction** consists of pointwise adjoint bifunctors

$$\mathcal{K} \times \mathcal{M} \xrightarrow{\times} \mathcal{N} \quad \mathcal{K}^{\text{op}} \times \mathcal{N} \xrightarrow{\text{hom}_\ell} \mathcal{M} \quad \mathcal{M}^{\text{op}} \times \mathcal{N} \xrightarrow{\text{hom}_r} \mathcal{K}$$

$$\mathcal{N}(k \times m, n) \cong \mathcal{M}(m, \text{hom}_\ell(k, n)) \cong \mathcal{K}(k, \text{hom}_r(m, n))$$

## Examples

A closed monoidal structure  $(\times, \text{hom}_\ell, \text{hom}_r): \mathcal{V} \times \mathcal{V} \rightarrow \mathcal{V}$ .

A tensored and cotensored enriched category  $(\otimes, \{\}, \text{hom}): \mathcal{V} \times \mathcal{M} \rightarrow \mathcal{M}$ .

# The monoidal model structure on **sSet**

To prove that **sSet** is a **monoidal model category**, suffices to show:

- $(\partial\Delta^n \rightarrow \Delta^n) \hat{\times} (\partial\Delta^m \rightarrow \Delta^m)$  is a cofibration
- $(\Lambda_k^n \rightarrow \Delta^n) \hat{\times} (\partial\Delta^m \rightarrow \Delta^m)$  is an anodyne extension
- $(\partial\Delta^m \rightarrow \Delta^m) \hat{\times} (\Lambda_k^n \rightarrow \Delta^n)$  is an anodyne extension

Analogously, though this was very hard to prove:

## Cellularity & Uniqueness Theorem (R.)

A cofibrantly generated algebraic model structure on a closed monoidal category is a **monoidal algebraic model structure** if and only if the pushout-products of the generating (trivial) cofibrations are cellular. The assignment of coalgebra structures to these maps completely determines the constituent **algebraic Quillen two-variable adjunction**.

## Corollaries (R.)

**sSet** and **Cat** are monoidal algebraic model categories

# A peek behind the curtain

The proof uses a composition criterion:

## Theorem (R.)

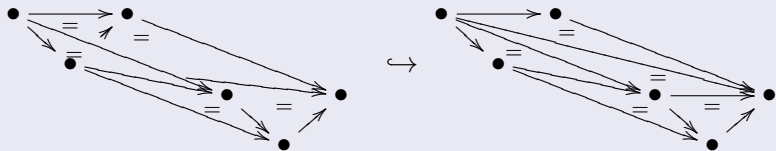
A lifted functor  $\hat{\text{hom}}(-, -)$  determines a **two-variable adjunction of awfs** iff, given a coalgebra  $i$  and composable algebras  $f, g$ , the algebra  $\hat{\text{hom}}(i, gf)$  solves a lifting problem against a coalgebra  $j$  as follows:

$$\begin{array}{ccccccc}
 K & \xrightarrow{a} & X^{B^a} & \xrightarrow{\cong} & X^B & \xrightarrow{f^B} & Y^B \\
 \downarrow j & \nearrow e & \downarrow \hat{\text{hom}}(i,f) & \nearrow e & \downarrow \hat{\text{hom}}(i,gf) & \nearrow d & \downarrow \hat{\text{hom}}(i,g) \\
 L & \xrightarrow{d \times c} & Y^B \times_{Y^A} X^A & \xrightarrow{g^B \times_{g^A} 1} & Z^B \times_{Z^A} X^A & \xrightarrow{1 \times_1 f^A} & Z^B \times_{Z^A} Y^A \\
 & \searrow b \times c & & & & & 
 \end{array}$$

and also satisfies a dual condition in the first variable.

# A combinatorial tidbit

The monoidal algebraic model structure on **sSet** defines (a priori) two different “anodyne cell structures” for the pushout-product of two anodyne cell complexes—eg,  $(\Lambda_1^2 \rightarrow \Delta^2) \hat{\times} (\Lambda_0^1 \rightarrow \Delta^1)$



—and these *are* different:

- one fills the missing end triangle and then the “trough”
- the other fills in the top square and then the interior cylinder

Future work will explore the implications of these results for the theory of enriched model categories.

# Acknowledgments

## Thanks

Thanks to the organizers, Richard Garner, Martin Hyland, Peter May, Mike Shulman, and the members of the category theory seminars at Chicago, Macquarie, and Sheffield.

## For further details

- Riehl, E., Algebraic model structures, *New York J. Math* 17 (2011) 173-231.
- Riehl, E., Monoidal algebraic model structures, arXiv:1109.2883v1 [math.CT], or at [www.math.harvard.edu/~eriehl](http://www.math.harvard.edu/~eriehl)